

## Lena Funcke

## Illit <br> $\therefore C^{2} Q A$ Co-design Center for Quantum Advantage


with Karl Jansen, Stefan Kühn, Tobias Hartung, et al.
Lattice 2022 Conference, Bonn, 8 August 2022

## Why quantum computing?

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Computational costs of lattice field theory
Supercomputer usage for different fields (INCITE 2019)

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$\rightarrow$ Lattice QCD: ~ 40\%

Figure credit:
Jack Wells, Kate Clark


- Astrophysics
- Al-Materials
- Nuclear Physics
- Biophysics
- Plasma Physics
- Seismology

- Subsurface Flow . Weather/Climate


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Computational challenges of lattice field theory
Critical slowing down, large autocorrelation times, ...
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## Why is the sign problem exponentially hard?

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## Partition function <br> $Z=\int D U D \bar{\psi} D \psi e^{-S}=\int D U e^{-S_{g}} \operatorname{det} M$

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## Example: finite baryon chemical potential

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Importance sampling
Interpretation of $e^{-S_{g}} \operatorname{det} M$ as probability weight

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## Reweighting procedure

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Example: phase quenched theory
$\langle O\rangle=\frac{\int D U e^{-S_{g}}|\operatorname{det} M| e^{i \phi} O}{\int D U e^{-S_{g}}|\operatorname{det} M| e^{i \phi}}=\frac{\left\langle e^{i \phi} O\right\rangle_{p q}}{\left\langle e^{i \phi}\right\rangle_{p q}}$

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Baryon density
Figure credit:
BNL/RHIC,
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Highly oscillatory integrands
Near-cancellation of positive \& negative contributions


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\int d x \exp \left(-x^{2}+i \lambda x\right) \rightarrow \int d x \exp \left(-x^{2}\right) \cos (\lambda x)
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## Highly oscillatory integrands

Near-cancellation of positive \& negative contributions
Sample number grows exponentially with volume $V$


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## Classical approaches to tackle the sign problem

Why is(n't) classical computing enough?

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Orus (2014)


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Deep learning for path integral contour deformations ${ }^{1}$

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Nakayama et al. (2022)


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Simulate chemical potential, $\theta$-term, real-time dynamics ${ }^{2}$ Mostly focus on 1+1D, first simulations in 2+1D \& 3+1D3


Nakayama et al. (2022) Bañuls et al. (2013) Challenges

No efficient parametrization of highly entangled states
In real-time evolution, tensor size can grow exponentially

## Do we really need quantum computing?

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Example: 1+1D Bose-Hubbard model
Hamiltonian
$\mathcal{H}=\sum_{j}-J\left(\hat{a}_{j}^{\dagger} \hat{a}_{j+1}+\right.$ h.c. $)+\frac{U}{2} \hat{n}_{j}\left(\hat{n}_{j}-1\right)$


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## Practical quantum advantage in quantum simulation ${ }^{2}$

Here we overview the state of the art and future perspectives for quantum simulation, arguing that a first practical quantum advantage already exists in the case of specialized applications of analogue devices

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## Quantum computing: where do we stand?

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## Quantum hardware

## Achievements

Quantum advantage: outperformed classical computers ${ }^{1}$


Arute et al. (2019)


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Astibuag (2022)
$\left.\begin{array}{l}|1\rangle=8 \\ |0\rangle=\end{array}\right\} \quad|\psi\rangle=$

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$\mathcal{O}(100)$ digital / $\mathcal{O}(1000)$ analog qubits $\rightarrow$ need more


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## Applications

cryptography, optimization problems, ...
particle / nuclear / condensed matter physics, ...

## Challenges

new technology $\rightarrow$ need fundamentally new algorithms competition $\rightarrow$ classical algorithms quickly advance


## Quantum computing: where do we stand?

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Example: "quantum advantage" (2019)
Example: "classical advantage" (2021)

## Quantum supremacy using a programmable superconducting processor

The promise of quantum computers is that certain computational tasks might be executed exponentially faster on a quantum processor than on a classical processor ${ }^{1}$.A fundamental challenge is to build a high-fidelity processor capable of running quantum algorithms in an exponentially large computational space. Here we report the use of a processor with programmable superconducting qubits ${ }^{2-7}$ to create quantumstates on 53 qubits, corresponding to a computational state-space of dimension $2^{53}$ (about $10^{16}$ ). Measurements from repeated experiments sample the resulting probability distribution, which we verify using classical simulations. Our Sycamore processor takes about 200 seconds to sample one instance of a quantum circuit a million times-our benchmarks currently indicate that the equivalent task for a state-of-the-art classical supercomputer would take approximately 10,000 years. This dramatic increase in speed compared to all known classical algorithms is an experimental realization of quantum supremacy ${ }^{s-14}$ for this specific computational task, heralding a muchanticipated computing paradigm.

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Closing the "Quantum Supremacy" Gap: Achieving Real-Time Simulation of a Random Quantum Circuit Using a New Sunway Supercomputer

We develop a high-performance tensor-based simulator for random quantum circuits(RQCs) on the new Sunway supercomputer. Our major innovations include: (1) a near-optimal slicing scheme, and a path-optimization strategy that considers both complexity and compute density; (2) a threelevel parallelization scheme that scales to about 42 million cores; (3) a fused permutation and multiplication design that improves the compute efficiency for a wide range of tensor contraction scenarios; and (4) a mixed-precision scheme to further improve the performance. Our simulator effectively expands the scope of simulatable RQCs to include the $10 \times 10$ (qubits) $\times(1+40+1)$ (depth) circuit, with a sustained performance of 1.2 Eflops (single-precision), or 4.4 Eflops (mixed-precision)as a new milestone for classical simulation of quantum circuits; and reduces the simulation sampling time of Google Sycamore to 304 seconds, from the previously claimed 10,000 years.

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53- and 54-Qubit Sycamore Circuits with Single
Precision Storage to Disk (8 bytes per amplitude)


Pednault et al. (2019) runtime improved from 2.5 days (2019) to 304 seconds (2021)

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## $\rightarrow$ Quantum-classical race:

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The promise of quantum computers is that certain computational tasks might be executed exponentially faster on a quantum processor than on a classical processor ${ }^{1}$.A fundamental challenge is to build a high-fidelity processor capable of running quantum algorithms in an exponentially large computational space. Here we report the use of a processor with programmable superconducting qubits ${ }^{2-7}$ to create quantumstates on 53 qubits, corresponding to a computational state-space of dimension $2^{53}$ (about $10^{16}$ ). Measurements from repeated experiments sample the resulting probability distribution, which we verify using classical simulations. Our Sycamore processor takes about 200 seconds to sample one instance of a quantum circuit a million times-our benchmarks currently indicate that the equivalent task for a state-of-the-art classical supercomputer would take approximately 10,000 years. This dramatic increase in speed compared to all known classical algorithms is an experimental realization of quantum supremacy ${ }^{s-14}$ for this specific computational task, heralding a muchanticipated computing paradigm.

## $\rightarrow$ Quantum-classical race:

algorithms and hardware quickly advance
$\rightarrow$ For exponentially hard problems:
small quantum step $\leftrightarrow$ giant classical leap

Example: "classical advantage" (2021)

Closing the "Quantum Supremacy" Gap: Achieving Real-Time Simulation of a Random Quantum Circuit Using a New Sunway Supercomputer

We develop a high-performance tensor-based simulator for random quantum circuits(RQCs) on the new Sunway supercomputer. Our major innovations include: (1) a near-optimal slicing scheme, and a path-optimization strategy that considers both complexity and compute density; (2) a threelevel parallelization scheme that scales to about 42 million cores; (3) a fused permutation and multiplication design that improves the compute efficiency for a wide range of tensor contraction scenarios; and (4) a mixed-precision scheme to further improve the performance. Our simulator effectively expands the scope of simulatable RQCs to include the $10 \times 10$ (qubits) $\times(1+40+1)$ (depth) circuit, with a sustained performance of 1.2 Eflops (single-precision), or 4.4 Eflops (mixed-precision)as a new milestone for classical simulation of quantum circuits; and reduces the simulation sampling time of Google Sycamore to 304 seconds, from the previously claimed 10,000 years.

53- and 54-Qubit Sycamore Circuits with Single
Precision Storage to Disk (8 bytes per amplitude)


Note: classical runtime improved from 2.5 days (2019) to 304 seconds (2021)

Pednault et al. (2019)

## Quantum computing: where will we go?

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The Path to Go...
A Rough Sketch...
State of the Art
IBM: 27 physical qubits $(2019) \rightarrow 65(2020) \rightarrow 127(2021)$

Carlow (2018)

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Need \(\mathcal{O}\left(10^{7}-10^{8}\right)\) logical qubits for lattice volume of \(96^{3}\)

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A Rough Sketch...


Gidney, Ekera (2019); Kan, Nam (2021)

\section*{How can we reduce the noise?}

Noisy quantum circuit

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\section*{Error mitigation versus error correction}

\section*{Problem}

Quantum noise: affecting qubits, gates, measurement

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Error correction (EC): fault-tolerant quantum computation

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Error correction (EC): fault-tolerant quantum computation E.g. bit-flip code, \({ }^{1}\) Shor code, \({ }^{2}\) toric code, \({ }^{3}\) GKP code, \({ }^{4} \ldots\)

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For EC, need extra qubits and noise below threshold \({ }^{5}\)
E.g. surface code needs \(>1000\) extra qubits for \(p<0.1 \%\)

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Operator rescaling method
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mitigate bit-flip errors during readout: \(0 \xrightarrow{p_{0}} 1\) or \(1 \xrightarrow{p_{1}} 0\)

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\hline Readout & Bit Flips & Probability & Noisy Operator \\
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\hline & & \\
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\hline\(\ldots\) for outcome 0 & \(0 \rightarrow 1,1 \rightarrow 1\) & \(p_{0}\left(1-p_{1}\right)\) & \(\tilde{O}=-\mathbb{I}=\left(\begin{array}{cc}-1 & 0 \\
0 & -1\end{array}\right)\) \\
\hline\(\ldots\) for outcome 1 & \(0 \rightarrow 0,1 \rightarrow 0\) & \(\left(1-p_{0}\right) p_{1}\) & \(\tilde{O}=\mathbb{I}=\left(\begin{array}{ll}1 & 0 \\
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\hline
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\hline \begin{array}{c}\text { incorrect } \\
\ldots \text { for both outcomes }\end{array} & 0 \rightarrow 1,1 \rightarrow 0 & p_{0} p_{1} & \tilde{O}=-Z=\left(\begin{array}{cc}-1 & 0 \\
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\hline \ldots \text { for outcome } 0 & 0 \rightarrow 1,1 \rightarrow 1 & p_{0}\left(1-p_{1}\right) & \tilde{O}=-\mathbb{I}=\left(\begin{array}{cc}-1 & 0 \\
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\hline \ldots \text { for outcome } 1 & 0 \rightarrow 0,1 \rightarrow 0 & \left(1-p_{0}\right) p_{1} & \tilde{O}=\mathbb{I}=\left(\begin{array}{cc}1 & 0 \\
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\(\left.=\begin{array}{c}\text { Total noisy operator: } \tilde{O} \\
+p_{0}\left(1-p_{1}\right)\left(1-p_{1}\right) Z+p_{0} p_{1}(-Z)+\left(1-p_{0}\right) p_{1} \mathbb{I}\end{array}\right]\)
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\begin{aligned}
& =\left(1-p_{0}\right)\left(1-p_{1}\right) \boldsymbol{Z}+p_{0} p_{1}(-Z) \\
& +p_{0}\left(1-p_{1}\right)(-\mathbb{I})+\left(1-p_{0}\right) p_{1} \mathbb{I}
\end{aligned}
\] \\
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\hline ... for outcome 1 & \(0 \rightarrow 0,1 \rightarrow 0\) & \(\left(1-p_{0}\right) p_{1}\) & \(\tilde{O}=\mathbb{I}=\left(\begin{array}{ll}1 & 0 \\ 0 & 1\end{array}\right)\) & \(\rightarrow \boldsymbol{Z}=\frac{1}{1-p_{0}-p_{1}} \tilde{O}-\frac{p_{1}-p_{0}}{1-p_{0}-p_{1}} \mathbb{I}\) \\
\hline
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Example: gate error mitigation

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Benchmark: \(Z\) and \(Z_{1} Z_{2}\) operators on IBM-Q hardware


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Benchmark: \(Z\) and \(Z_{1} Z_{2}\) operators on IBM-Q hardware
Result: measurement error reduced by factor 10



Carlow (2018)

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\section*{Other mitigation techniques}

Zero-noise extrapolation, \({ }^{2}\) randomized compiling, \({ }^{3}\) quasi-probability decomposition, \({ }^{4}\)...
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\section*{Lattice field theory applications}

Zero-noise extrapolation for lattice Schwinger model:


Klco et al. (2018)
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Experimental results on "public" QC
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Real-time evolution: Schwinger model, \({ }^{1} \mathrm{SU}(2),{ }^{2} \mathrm{SU}(3),{ }^{3}\)...
\(1+1 \mathrm{D} \operatorname{SU}(3)\) gauge theory, one plaquette \({ }^{3}\)


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Variational computation: SU(2) "hadron" masses \({ }^{4}\)

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Cold atoms
Analog simulation: Bose-Hubbard, \({ }^{6}\) Schwinger model, \({ }^{7}\)..

\section*{How to simulate these field theories?}

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\section*{Key concept}

Classical computer: main computation

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\section*{Goal}

Find ground state and excited states of Hamiltonian \(\mathcal{H}\)

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Minimize \(E(\vec{\alpha})=\langle\psi(\vec{\alpha})| \mathcal{H}|\psi(\vec{\alpha})\rangle\) w.r.t. parameters \(\vec{\alpha}\)

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Given \(E\left(\vec{\alpha}_{i}\right)\), find optimized parameters \(\vec{\alpha}_{i+1}\)


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Sim et al. (2018)

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Encode


Sim et al. (2018)

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\section*{How to prepare the quantum state?}

\section*{Quantum circuit design}

Example: geometrical method \({ }^{2}\)

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Gauge invariance requires imposing local constraints

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Summary: where do we stand, where will we go?
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The Path to Go...
A Rough Sketch...

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new and quickly progressing
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\section*{Classical}
more established but limited
\(\rightarrow\) solid, promising


A Rough Sketch...


Gidney, Ekera (2019); Kan, Nam (2021)

\section*{Summary: where do we stand, where will we go?}

The Way Forward...

\section*{Quantum}
new and quickly progressing
\(\rightarrow\) high risk, high gain


\section*{Classical}
more established but limited
\(\rightarrow\) solid, promising

\section*{Combination}
\(\rightarrow\) quantum-classical algorithms

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\section*{Backup: details of Bose-Hubbard quantum simulation}

\section*{Classical versus quantum simulation}

Why is(n't) classical computing enough?

\section*{Bose-Hubbard Hamiltonian}
\[
\mathcal{H}=\sum_{j}-J\left(\hat{a}_{j}^{\dagger} \hat{a}_{j+1}+\text { h.c. }\right)+\frac{U}{2} \hat{n}_{j}\left(\hat{n}_{j}-1\right)+\frac{K}{2} \hat{n}_{j} j^{2}
\]

\section*{Experimental goal \({ }^{1}\)}

Simulate quantum tunneling from even to odd sites


\section*{Experimental setup}

Quantum simulation: ultracold atoms in optical lattice Classical benchmark: tensor networks (MPS-based)

\footnotetext{
\({ }^{1}\) Trotzky et al. (2012)
}

\section*{Experimental results}

Classical (line) vs. quantum (circles) simulation

this dynamical quantum simulator outperforms any continuous-time numerical simulation, for which the calculational effort increases with \(t\). Simulation methods on classical computers, such as matrix-product state based \(t\)-DMRG used here, suffer from an extensive increase in entanglement entropy which limits the relaxation times accessible in the calculations \({ }^{24,25}\).

\section*{Backup: analog versus digital quantum computers}

\section*{Analog}

\section*{Digital}

\section*{Concept}
- Use controllable quantum system to simulate the behavior of another quantum system
- Continuous time evolution
- Usually non-universal


\section*{Concept}
- Construct set of logical gates onto qubits
- Discrete time evolution
- Usually universal


\section*{Backup: how to measure the energy in VQE?}

\section*{Example: massless Schwinger model}

\section*{Original Hamiltonian}
\(\mathcal{H}=-\frac{i}{2 a} \sum_{n=0}^{N-2}\left(\phi_{n}^{\dagger} e^{i \theta_{n}} \phi_{n+1}-\right.\) h. c. \()+\frac{a g^{2}}{2} \sum_{n=0}^{N-2} F_{n}^{2}\)
with \(\theta_{n}=-a q A_{n}^{1}, g F_{n}=E_{n},\left[\theta_{n}, L_{m}\right]=i \delta_{n m}, \theta_{n} \in[0,2 \pi]\)
Eliminate \(\boldsymbol{\theta}_{\boldsymbol{n}}\)
\(\phi_{n}^{\dagger} e^{i \theta_{n}} \phi_{n+1} \rightarrow \phi_{n}^{\dagger} \phi_{n+1}\) from gauge transformation: \(\phi_{n} \rightarrow\left(\prod_{k=0}^{n-1} e^{-i \theta_{n}}\right) \phi_{n}\) and \(\phi_{n}^{\dagger} \rightarrow \phi_{n}^{\dagger}\left(\prod_{k=0}^{n-1} e^{i \theta_{n-k}}\right)\)
Eliminate \(\boldsymbol{F}_{\boldsymbol{n}}\)
\(F_{n}=\sum_{k=0}^{n} Q_{k}\) from solving Gauß law (for OBC):
\(F_{n}-F_{n-1}=Q_{n} \forall n\), where \(Q_{n}=\phi_{n}^{\dagger} \phi_{n}-\frac{1}{2}\left[1-(-1)^{n}\right]\)

\section*{Mapping the model to qubits}

\section*{Dimensionless spin Hamiltonian \({ }^{1}\)}
\(\mathcal{H}=x \sum_{n=0}^{N-2}\left(\sigma_{n}^{+} \sigma_{n+1}^{-}+\sigma_{n}^{-} \sigma_{n+1}^{+}\right)+\frac{1}{2} \sum_{n=0}^{N-2}\left\{\sum_{k=0}^{n}\left[(-1)^{k}+\sigma_{k}^{z}\right]\right\}^{2}\)
from mapping \(\phi_{n}^{\dagger} \phi_{n+1} \rightarrow \sigma_{n}^{+} \sigma_{n+1}^{-}\)and \(\phi_{n}^{\dagger} \phi_{n} \rightarrow \frac{1}{2}\left(\sigma_{n}^{Z}+\mathbb{I}\right)\)

\section*{Quantum computer}

Measurement of \(\langle\psi| \boldsymbol{O}|\psi\rangle\) with \(\boldsymbol{O} \in\left\{\mathbb{I}, \sigma^{z}\right\}^{\otimes N}\)
\(\mathcal{H}=\sum_{k} h_{k} U_{k}^{*} \boldsymbol{O}_{\boldsymbol{k}} U_{k}\) with \(U_{k}^{*} \boldsymbol{O}_{\boldsymbol{k}} U_{k} \in\left\{\mathbb{I}, \sigma^{x}, \sigma^{y}, \sigma^{z}\right\}^{\otimes N}\)


Gokhale et al. (2020)

\section*{Backup: quantum volume}

\section*{Concept}

\section*{Motivation}

Number of noisy qubits: no good performance measure New performance measure
Measure capabilities and error rates of quantum device

\section*{IBM's definition}
\[
\log _{2} V_{Q}=\arg \max _{\mathrm{n} \leq N}\{\min [n, d(n)]\}
\]

\section*{Example}

Successfully run circuit of depth \(d=8\) on \(n=8\) qubits: quantum volume is \(V_{Q}=2^{8}=256 \rightarrow\) size of state space "Success"
Most likely outputs of the circuit are computed correctly \(67 \%\) of the time with a \(2 \sigma\) confidence interval

\section*{Timeline}

\section*{Last three years}

Early 2020: \(V_{Q}=32\) (IBM) for \(d=5, n=5\)
Early 2021: \(V_{Q}=512\) (Honeywell) for \(d=9, n=9\)
Early 2022: \(V_{Q}=4096\) (Quantinuum) for \(d=12, n=12\)


\section*{Backup: preparing for overcoming sign problems in 3+1D}

\section*{Goal}

Simulate phase transition at \(\theta=\pi\) and large \(g=\beta^{-1 / 2}\)

\section*{Analytical results}

Derivation of Hamiltonian lattice \(\theta\)-term:
\(\theta Q=-\frac{i g^{2} \theta}{8 \pi^{2} a} \sum_{n, i, j, k, b} \varepsilon_{i j k} \operatorname{Tr}\left[E_{n, i}^{b} \lambda^{b}\left(U_{n, j k}-U_{n, j k}^{\dagger}\right)\right]\)

\section*{Numerical results}

Unlike in QCD, transition in \(U(1)\) might be not of first order

\section*{Near-future outlook}

Larger-volume simulation with 3+1D tensor networks

\section*{Far-future outlook}

First quantum simulation of \(3+1 \mathrm{D} \theta\)-term
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[^0]:    ${ }^{1}$ Trotzky et al. (2012)

[^1]:    ${ }^{1}$ Trotzky et al. (2012)

[^2]:    ${ }^{1}$ Trotzky et al. (2012)

